

Foundation University

Rawalpindi Campus

Introduction to Database Systems – CSC - 221 APresentation by R.M.Hafeez Javed WWW.rmbjaved.com



O MY NAME IS.....

✓ I REMEMBER......

Objective of Today's Lecture

Relations and Relational Algebra

RDM – Two Major Strengths

Simplicity – Relation Form

Strong Mathematical Foundation

Relational Data Model

- Presented by E. F. Codd in 1970
- Ø Before Relational Data Model, two older data models were in use;
- Hierarchical
- O Network

Basics of RDM

 Mainly used for external, conceptual, and to some extent physical schema.

 Separation of conceptual and physical levels makes manipulation much easier, contrary to previous data models.

Basics of RDM

- The basic structure is relation.
- O Relations physically represented as tables.
- Oconsists of rows and columns
- Both entities and relationships are modeled using tables/relations.

Basics of RDM

Columns represent attributes and rows represent records.

 Rows, records and tuples all these terms are used interchangeably.

A Table/Relation

| stID | stName | clName | doB | Gender |
|------|---------|--------|---------|--------|
| S001 | Zofeen | BCSE | 12/6/92 | F |
| S002 | Rohaf | BCSE | 3/9/92 | Μ |
| S003 | Noshail | BSCS | 7/8/93 | F |
| S004 | Rameen | BCA | 23/4/94 | F |
| S005 | Bilaval | BBA | 22/7/93 | Μ |

Relation - Terminology

- *Degree of a relation*: how long the tuples are, or how many columns the table has.
 - In the exemplary relation, degree of the relation is 5

- Cardinality of the relation: how many different tuples are there, or how many different rows the table has.
 - In the exemplary relation, cardinality of the relation is 5

Mathematical Relations

O Consider two sets

A = {x, y}
B = {2, 4, 6}

Cartesian product of these sets
 A X B = {(x,2), (x,4), (x,6), (y,2), (y,4), (y,6)}

Mathematical Relations

 A relation is some subset of this Cartesian product, For example

 $O R1 = \{(x,2), (y,2), (x,6), (x,4)\}$

 $O R2 = \{(x,4), (y,6), (y,4)\}$

Relational Model: Data Manipulation

- Data is represented as relations.
- Manipulation of data (guery and update operations) corresponds to operations on relations.
- Relational algebra describes those operations. They take relations as arguments and produce new relations.
- O Think of numbers and corresponding operators +,-,\, * or booleans and corresponding operators &,|,! (and, or, not).
- Relational algebra contains two kinds of operators: common set-theoretic ones and operators specific to relations (for example projecting on one of the columns).

Union

Standard set-theoretic definition of union:

 $\mathcal{O} \ \mathsf{A} \cup \mathsf{B} = \{ \mathsf{x} \colon \mathsf{x} \in \mathsf{A} \text{ or } \mathsf{x} \in \mathsf{B} \}$

✓ For example, {a,b,c} ∪ {a,d,e} = {a,b,c,d,e}

So we require in order to take a union of relations R and S that R and S have the same number of columns and that corresponding columns have the same domains.

Union – Compatible Relations

Two relations R and S are union-compatible if they have the same number of columns and corresponding columns have the same domains.

Example: Not Union Compatible

O Not Compatible!

O Different number of columns!

| Anne | aaa | 111111 | Tom | 1980 |
|-------|-----|--------|-------|------|
| Bob | bbb | 222222 | Sam | 1985 |
| Chris | CCC | 333333 | Steve | 1986 |

Example: Not Union Compatible

O Not Compatible!

O Different domains for the second column!

| Anne | aaa | Tom | 1980 |
|-------|-----|-------|------|
| Bob | bbb | Sam | 1985 |
| Chris | CCC | Steve | 1986 |

Example: Union-Compatible

| Anne | 1970 |
|-------|------|
| Bob | 1971 |
| Chris | 1972 |

| Tom | 1980 |
|-------|------|
| Sam | 1985 |
| Steve | 1986 |

Union of Two Relations

✓ Let R and S be two union-compatible relations. Then their union R ∪ S is a relation which contains tuples from both relations:

 $\mathsf{R} \cup \mathsf{S} = \{ x : x \in \mathsf{R} \text{ or } x \in \mathsf{S} \}.$

 Note that union is a partial operation on relations: it is only defined for some (compatible) relations, not for all of them.

Similar to division for numbers (result of division by 0 is not defined).

Example: Shopping Lists

| R | | S | | R U | S |
|--------|------|-------|------|------------|------|
| Cheese | 1.34 | Cream | 5.00 | Cheese | 1.34 |
| Milk | 0.80 | Soap | 1.00 | Milk | 0.80 |
| Bread | 0.60 | | | Bread | 0.60 |
| Eggs | 1.20 | | | Eggs | 1.20 |
| Soap | 1.00 | | | Soap | 1.00 |
| | | | | Cream | 5 00 |

5.00

Difference of Two Relations

Let R and S be two union-compatible relations. Then their *difference* R – S is a relation which contains tuples which are in R but not in S:

 $R - S = \{x: x \in R \text{ and } x \notin S\}.$

 Note that difference is also a partial operation on relations.

Example

| R | | S | | | R – S | | |
|--------|------|---------------|------|--|----------------|------|--|
| Cheese | 1 24 | Croom | E 00 | | Chaosa | 1 24 | |
| Milk | 0.80 | Cream Soap | 1.00 | | Cheese Milk | 0.80 | |
| Bread | 0.60 | | 1100 | | Bread | 0.60 | |
| Eggs | 1.20 | | | | Eggs | 1.20 | |
| Soap | 1.00 | | | | | | |

Intersection of Two Relations

✓ Let R and S be two union-compatible relations. Then their *intersection* is a relation R ∩ S which contains tuples which are both in R and S:

 $\mathsf{R} \cap \mathsf{S} = \{ x : x \in \mathsf{R} \text{ and } x \in \mathsf{S} \}$

 Note that intersection is also a partial operation on relations.

Intersection

| R | | S | | R \cap |
|--------|------|-------|------|-----------------|
| Cheese | 1.34 | Cream | 5.00 | Soap |
| Milk | 0.80 | Soap | 1.00 | |
| Bread | 0.60 | | | |
| Eggs | 1.20 | | | |
| Soap | 1.00 | | | |
| | | | | |

S

1.00

CHHUTTI

AND THAT IS FAREWELL TO DAY 15 ©